#### **Congestion Avoidance and Control**

Van Jacobson

Lawrence Berkeley Laboratory

#### Michael J. Karels

University of California at Berkeley

November, 1988

### Outline

# Slow-start Round-trip time timing Congestion avoidance

#### Packet Conservation Fails

- The connection doesn't get into equilibrium
- A sender injects a new packet before an old packet has exited
- The equilibrium can't be reached because of resource limits along the path





Getting to Equilibrium: Slow-start

Add a congestion window, cwnd, to the per-connection state

- When starting or restarting after a loss, set cwnd to one packet
- On each ack for new data, increase cwnd by one packet
- Sending the minimum of the receiver's advertised window and cwnd





### Conservation at equilibrium: round-trip timing

Good round trip time estimator
 TCP protocol specification suggests

$$R \leftarrow \alpha R + (1 - \alpha)M, \ \alpha = 0.9$$

rto  $\leftarrow \beta R, \beta = 2$ 

Backoff after a retransmit



Window adjustment interaction with round-trip timing Bandwidth-dominated  $\bullet$ In worst case,  $R_{i+1} = 2R_i$ •  $V_i = R_i - R_{i-1} = R_i / 2$  $\blacksquare RTO_i = R_i + 4 V_i$  $= 3 R_{i}$  $> 2 R_i$  $> R_{i+1}$ 

11



Adapting to the path: congestion avoidance Signal the congestion Sender's policy Congestion avoidance • On any timeout, *cwnd=cwnd*/2 On each ack for new data, cwnd+=1/cwnd Sending the minimum of the receiver's advertised window and cwnd



#### Multiple, simultaneous TCP



w/o congestion avoidance

with congestion avoidance

15







#### Future work: gateway side of congestion control

## Ensure fair sharing of the capacity Using packet drops as a congestion signal