

# **Experience on RSVP QoS Router Implementation**

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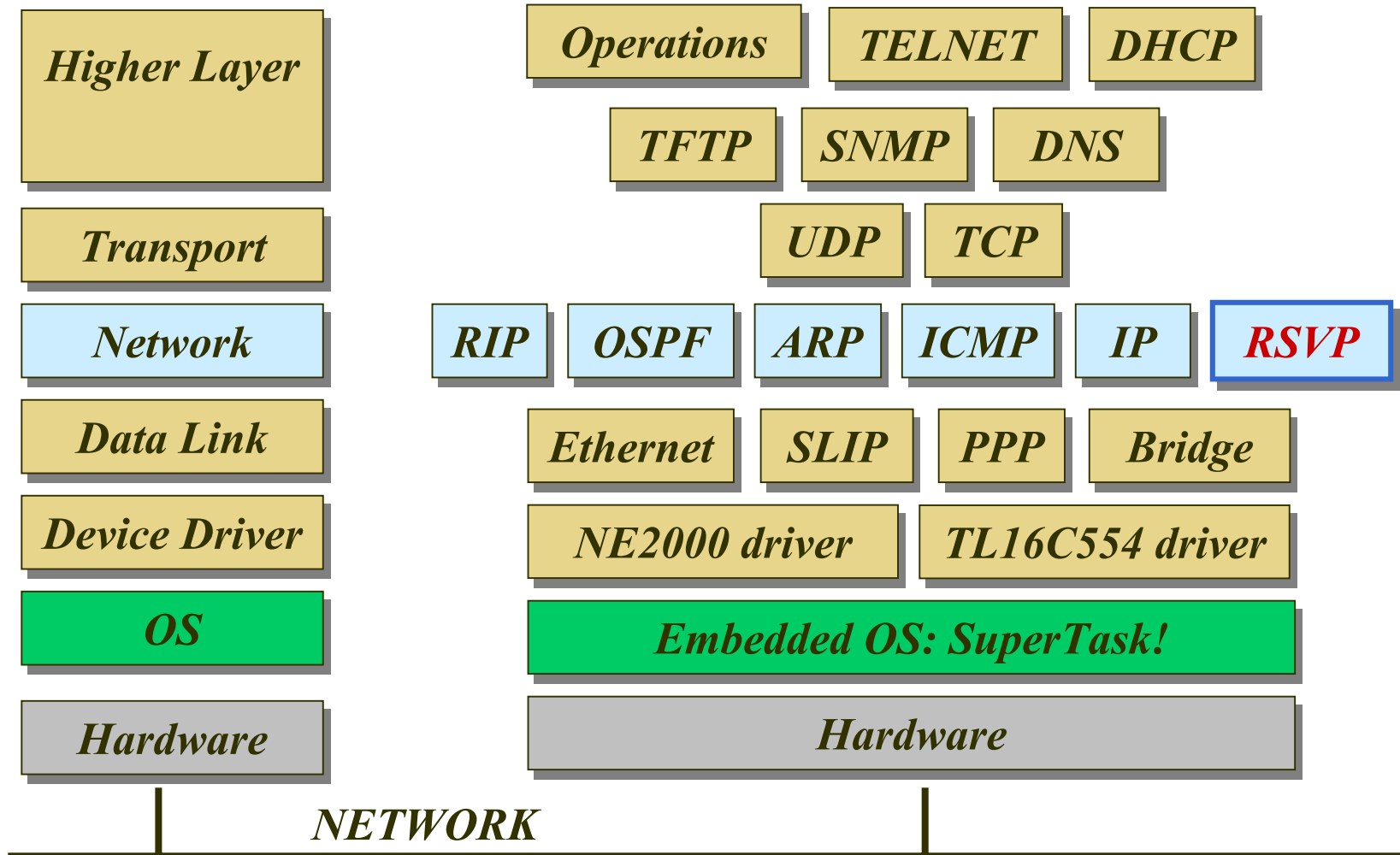
**National Chiao Tung University**

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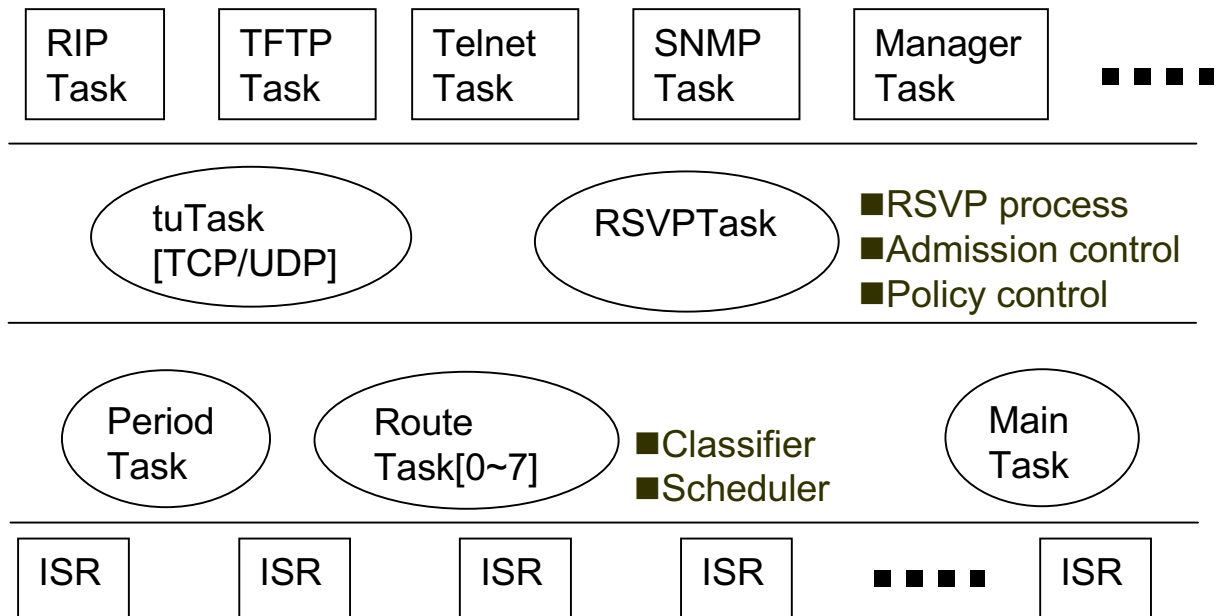
# Protocol Stack in RAS



Phase2 support
  USNET support

# Layered Tasks in RAS

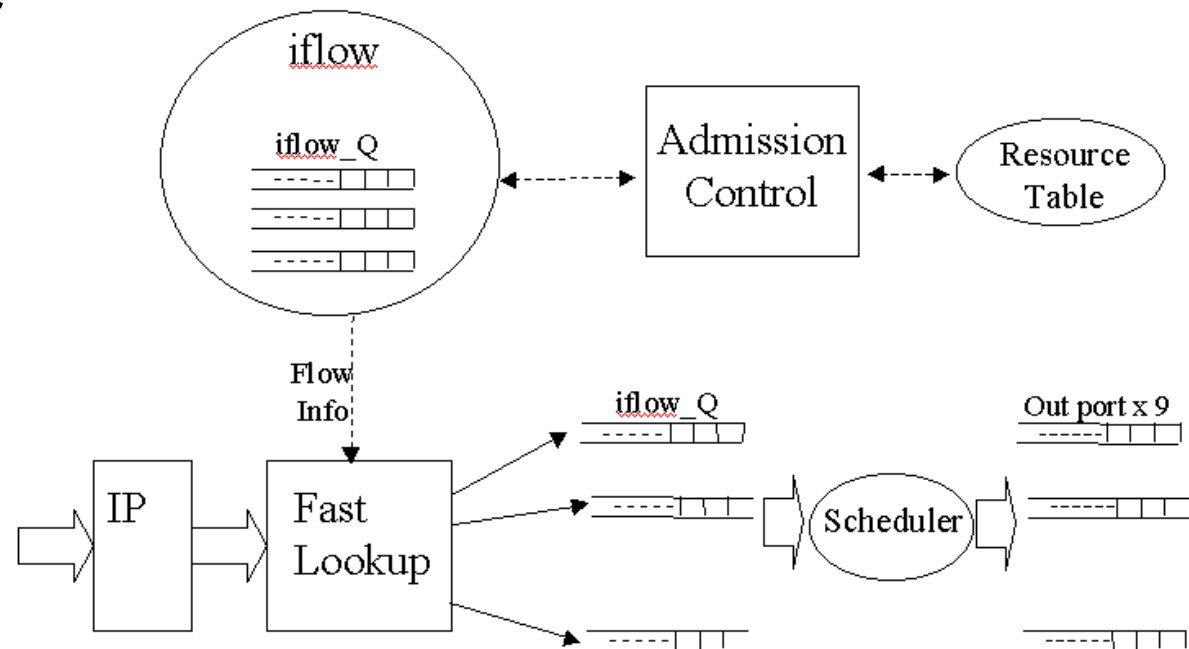
- The original CCL RAS (Remote Access Server) kernel just processes “**RSVP signaling**” messages.
- There are no “**traffic control modules**”.



# Traffic Control Modules

intercepting packets for QoS treatment

- Admission Control
- Classifier
- Scheduler



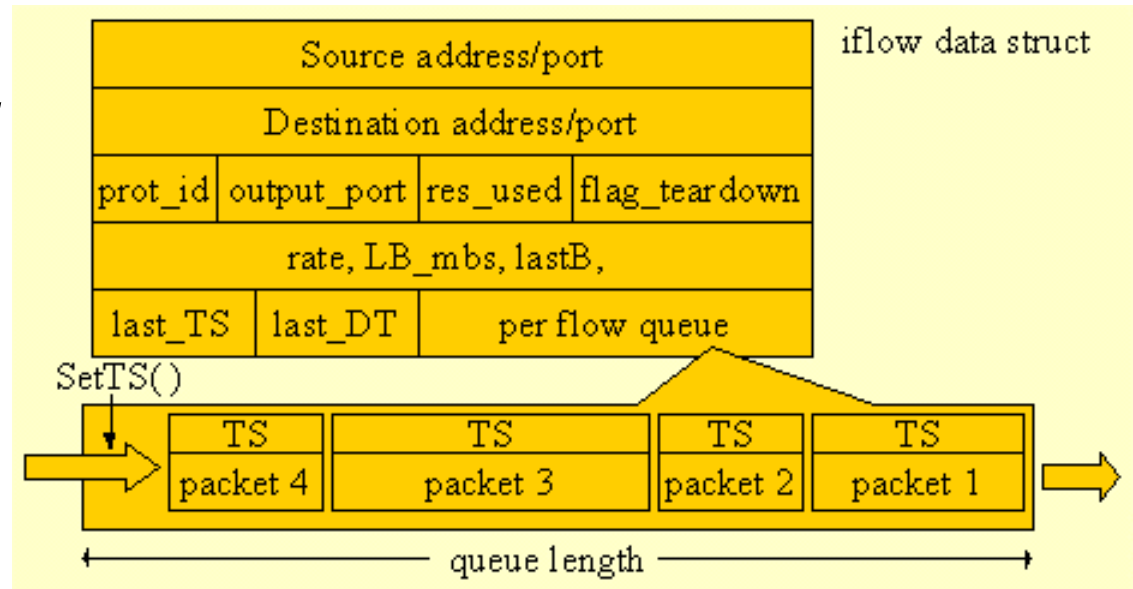
# Per-Flow Data Structures

each *FLOW*

→ *Classifier para.*

→ *Scheduler para.*

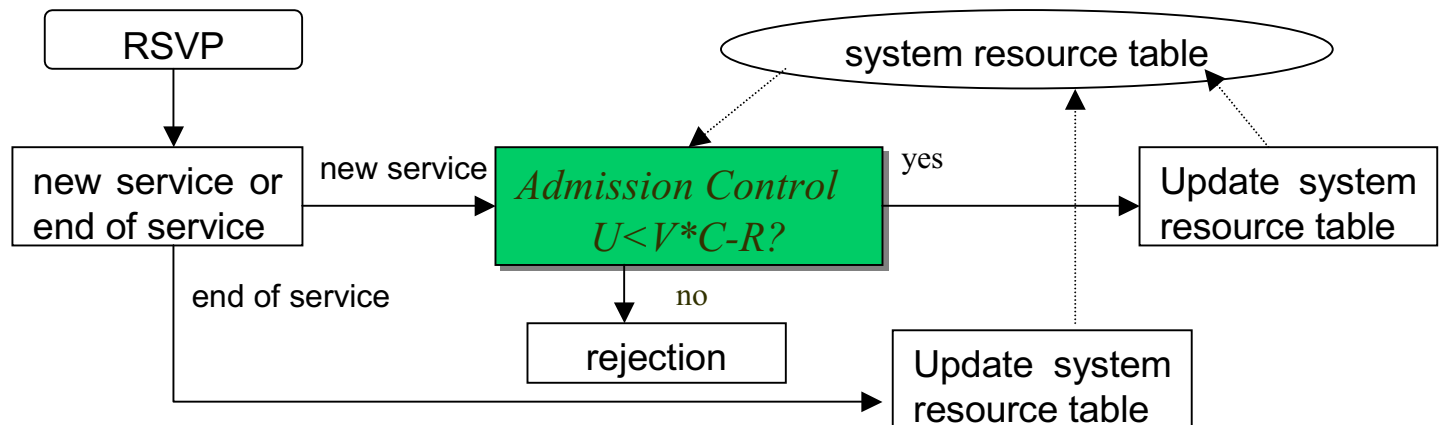
→ *Data queue*



# Process of Admission Control

Accept a new service request if ....  $U < V * C - R$

- Current utilization ( $U$ )
- Link capacity ( $C$ )
- Target utilization ( $V$ )
- Bandwidth of new request ( $R$ )



# System Resource Management

*After accepting a new request*

update data structure of iflow  
update **system resource table**

➤ **5 Flow IDs** for classifying  
➤ **Flow Spec.** for scheduling

*As the service ends*

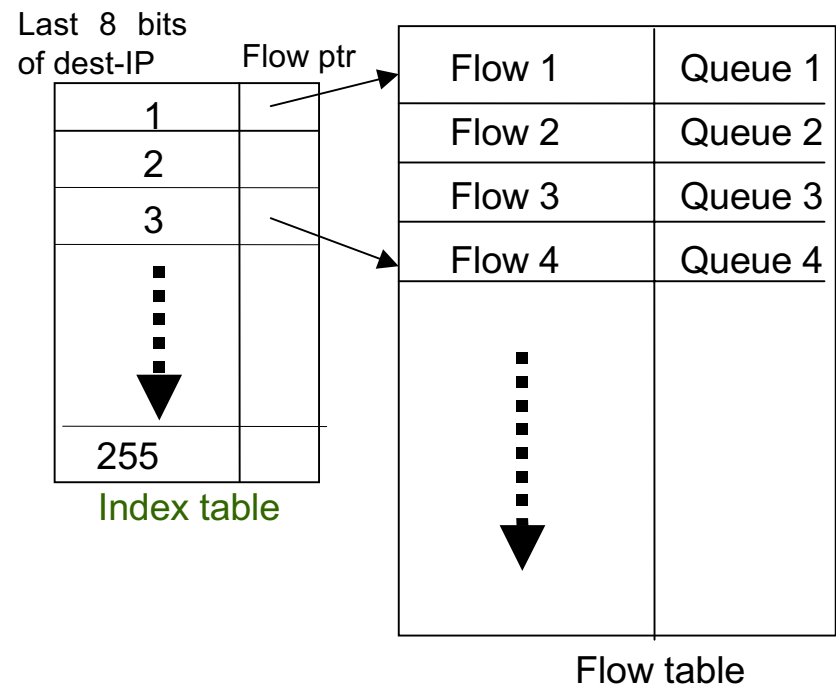
finish the transmission of those packets already in its queue  
update **system resource table**

port cid	0	1	2	..	8
C(bps)	10,000,000	33600	33600	..	33600
U(bps)				..	
v	0.9	0.9	0.9	..	0.9

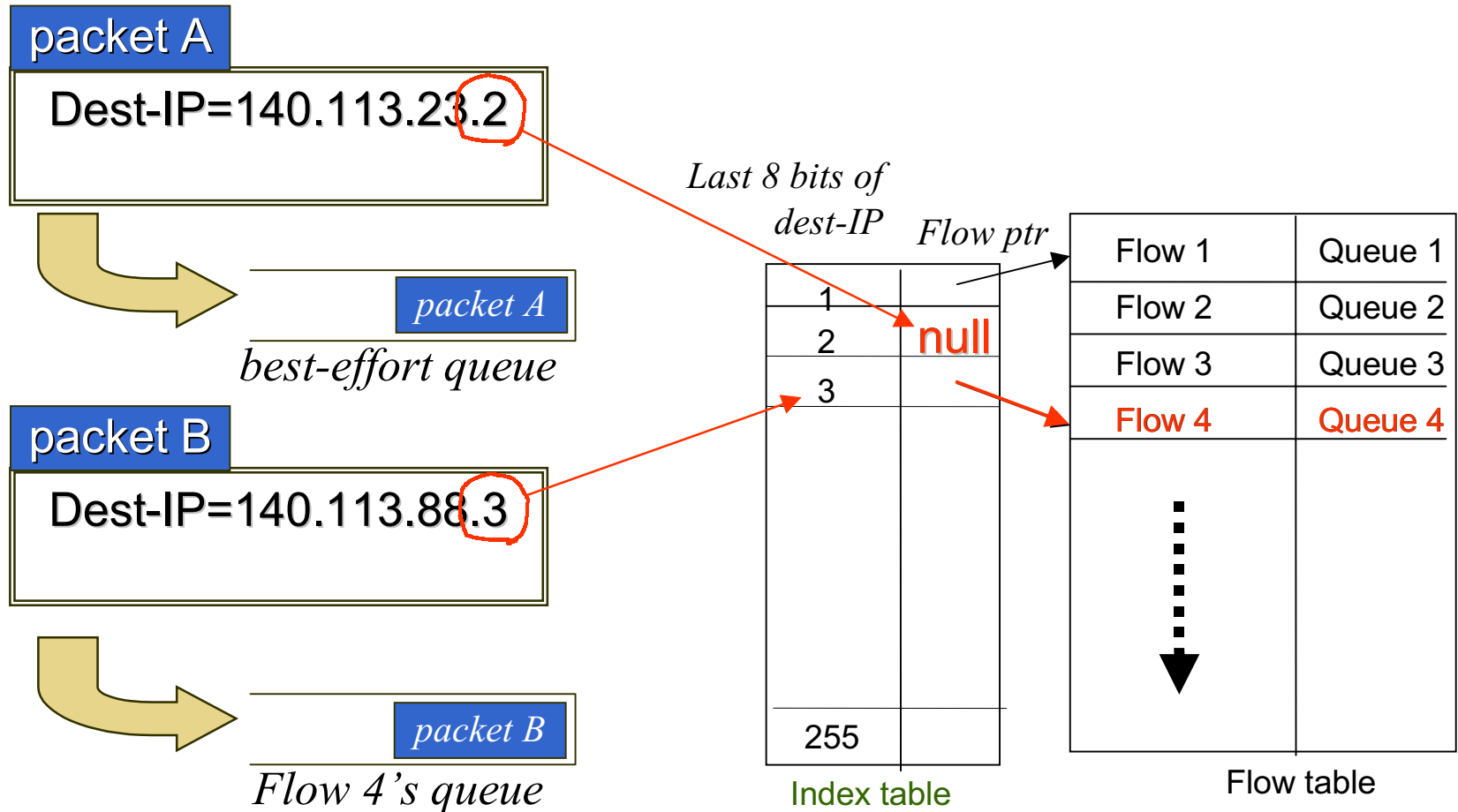


# Algorithm for Classifier

- 1 Look up the *last 8 bits of Dest-IP*
- 2 Get the flow's data and queue from **Flow Table** by its ptr in **Index Table**
- 3 Put the packet into the queue.



# Examples for Classifier



# Architecture of Scheduler

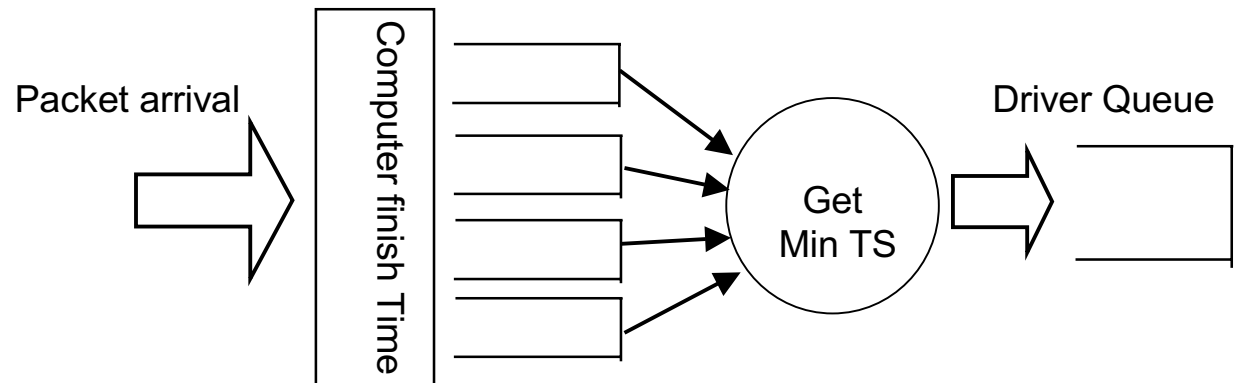
- ▶ Upon a packet arrival .....

Compute its finish time  $TS$

$$TS_i^k = \max(TS_i^{k-1}, v(t)) + \frac{L_i}{r_i}$$

- ▶  $v(t) = TS$  of the packet been serving as the  $k^{\text{th}}$  packet of flow  $i$  arrives (SCFQ)

- ▶ Scheduler continuously picks up the “**Min Finish Time**” packet among the first packets of all **queues**



# Conclusions

- ***A QoS-capable router needs:***
  - ***RSVP signaling for reservation***
  - ***traffic control modules to enforce the reservation***
  
- ***Scalability issues:***
  - ***number of flows***
  - ***number of packets***
  - ***bandwidth: link vs node***

# Reference

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